TOMORROW starts here.

Cisco Next Generation Fabric (Dynamic Fabric Automation Architecture)

BRKDCT-2385

Max Ardica – Technical Leader

Agenda

- DFA Requirements and Functions
- Fabric Management
- Workload Automation
- Optimized Network
  - Fabric Properties
  - Control Plane
  - Forwarding Plane
- Virtual Fabrics
- Hardware Support

Today’s DC Challenges

Are the result of…

Operational Complexity

Architecture Rigidity

Infrastructure Inefficiency

Dynamic

Fabric

Automation

Architecture

SIMPLIFY

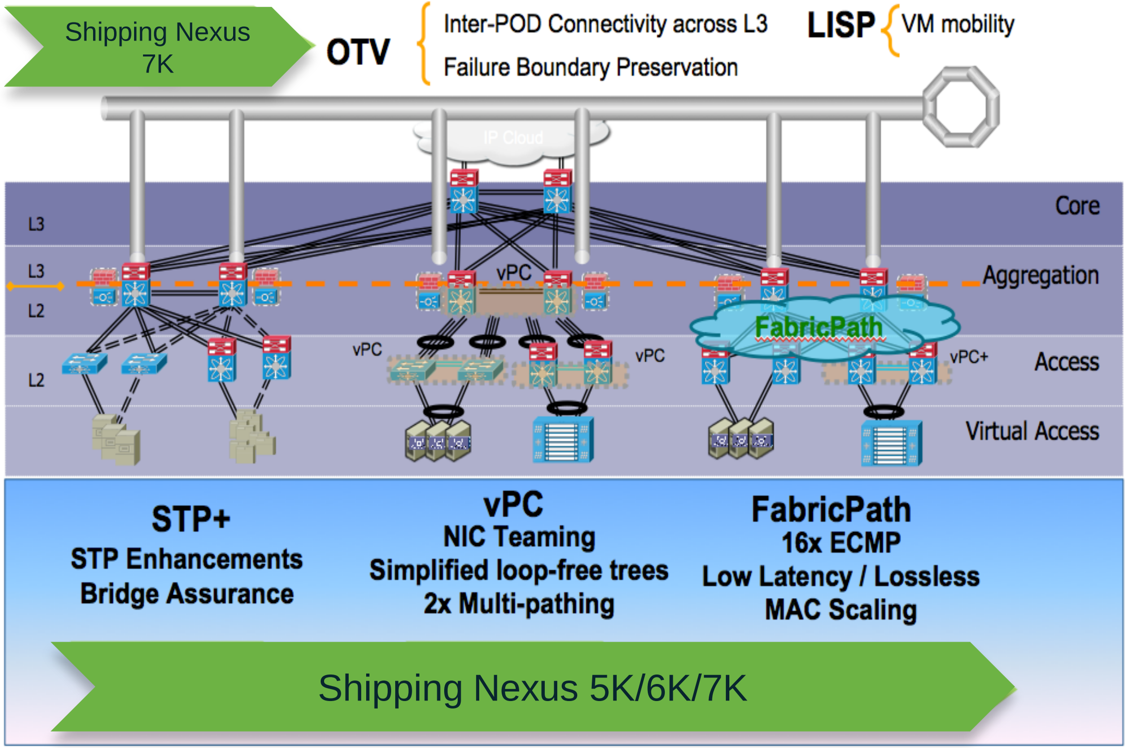
Auto-configuration at scale

OPTIMIZE

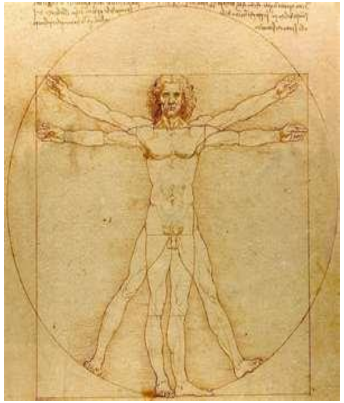
workload (physical/virtual) anywhere anytime

AUTOMATE

Dynamic Fabric Automation Architecture
The Latest Step in the POD Evolution Journey



… and now DFA!



Dynamic Fabric Automation Architecture

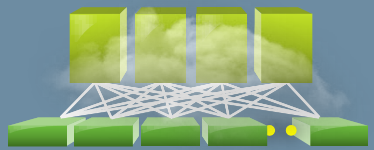
Innovative Building Blocks

Bundled functions are modular and simplified for scale and automation

Fabric Management

Workload Automation

Optimized Network



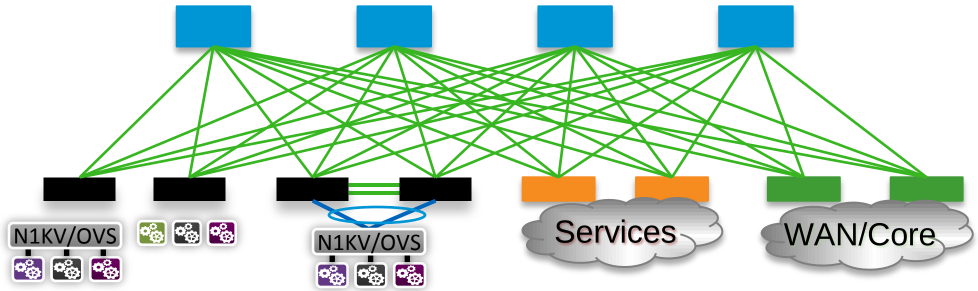
Virtual Fabrics

DFA Applicability and Use Cases

Cisco Dynamic Fabric Automation applies to any customer looking for solution to:

- ✓ DC Networks from the very small to the very large
- ✓ Environments with virtual and non-virtual workloads
- ✓ Looking to integrate with 3rd party Orchestration Tools
- ✓ Seeking flexibility on workload placement
- ✓ Looking for the Stability of small failure domains and flexibility or any app anywhere

Dynamic Fabric Automation Architecture
Device Roles



Spine
Leaf
Border Leaf
Services Leaf
Virtual Leaf* N1KV/OVS

Virtual Machines
Physical Machines
FEXs
3rd Party Switches
UCS FIs
Blade Switches
Storage

Firewalls
Load Balancers
3rd Party Appliances

Routers
Switches
3rd Party Devices

Note: the different leaf roles are logical and not physical. The same leaf switch could perform all three functions (regular, services and border leaf)

*Virtual Leaf: N1KV/OVS being a "light" participant on the control plane protocol (supporting VDP)

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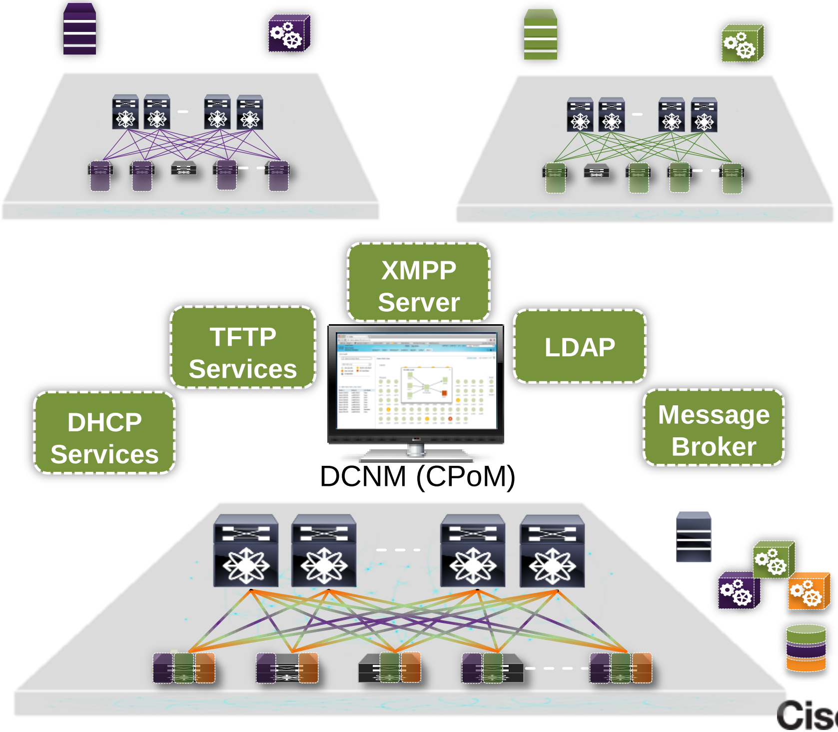
Simplifying Fabric Management & Optimizing Fabric Visibility

Fabric Management

Advantages

- Device Auto-Configuration
- Cabling Plan Consistency Check
- Automated Network Provisioning
- Common point of fabric access
- Network, vFabric & Host Visibility



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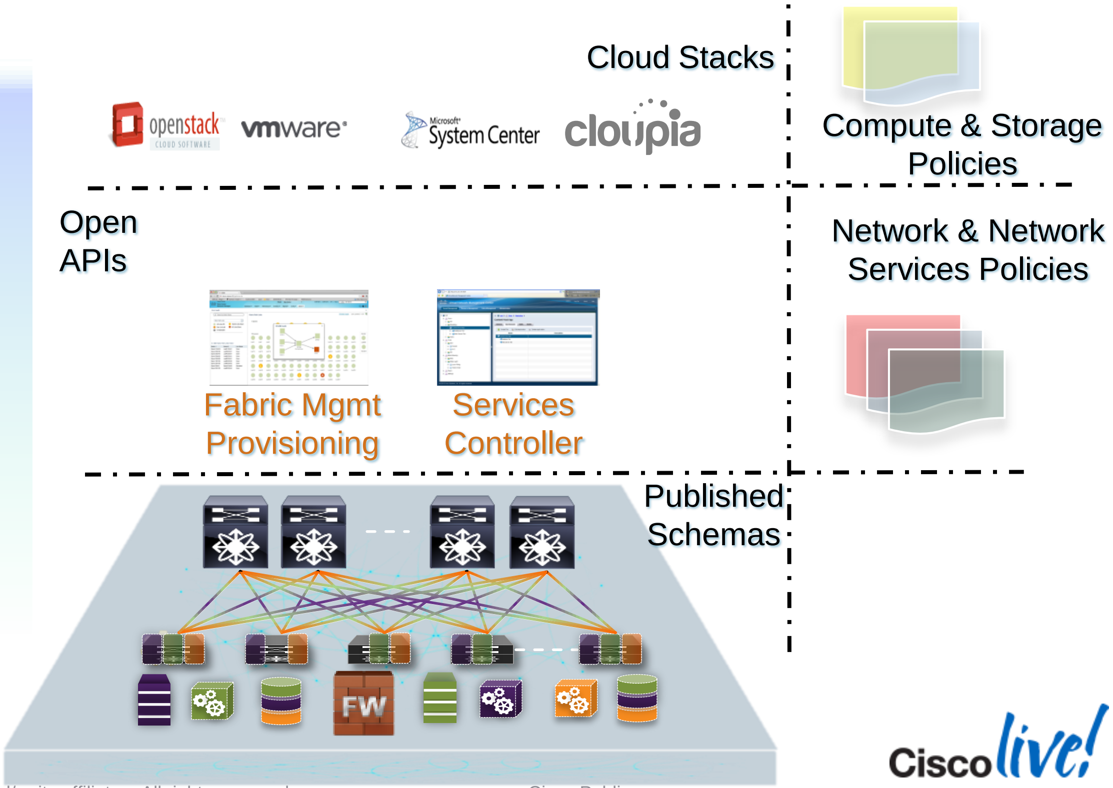
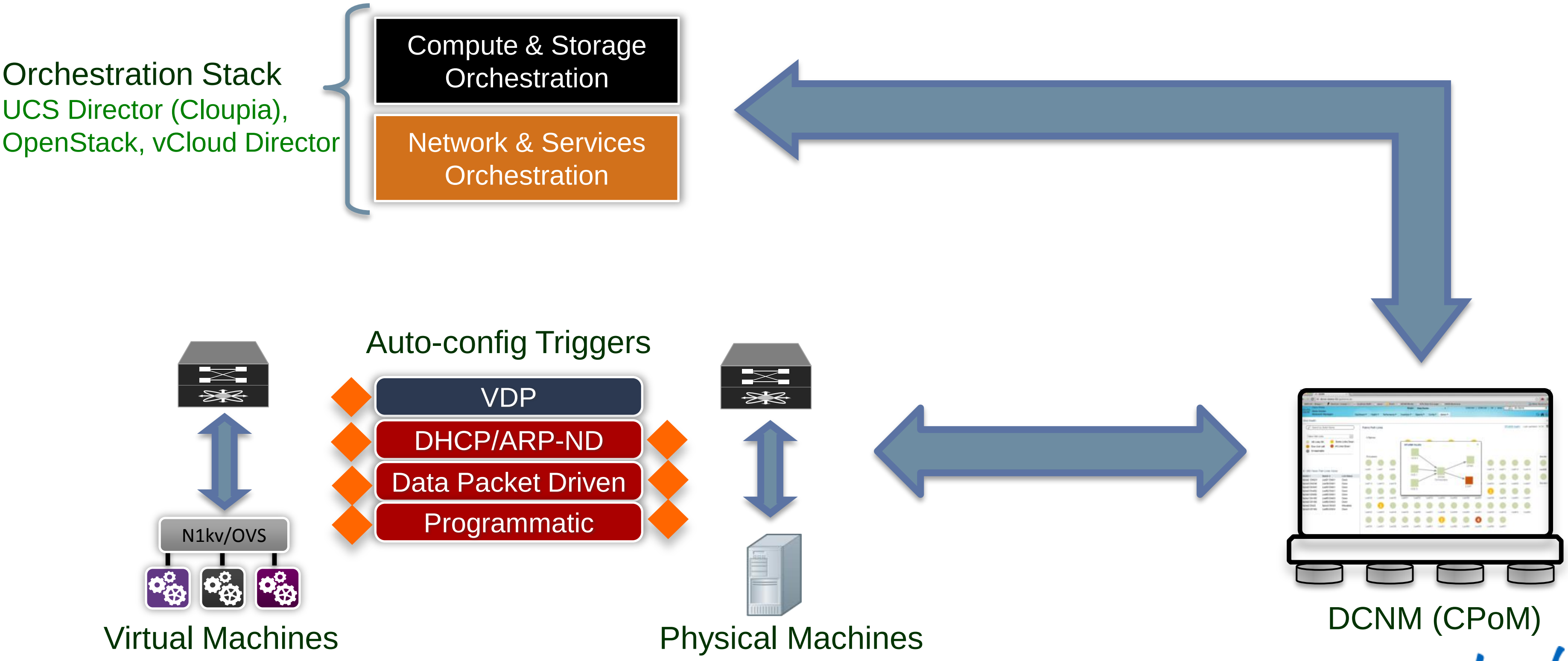
Workload Automation & Open Environment

Advantages

- Any workload, anywhere, anytime
- Open Integration: orchestration
- Automated scalable provisioning
- Workload aware fabric

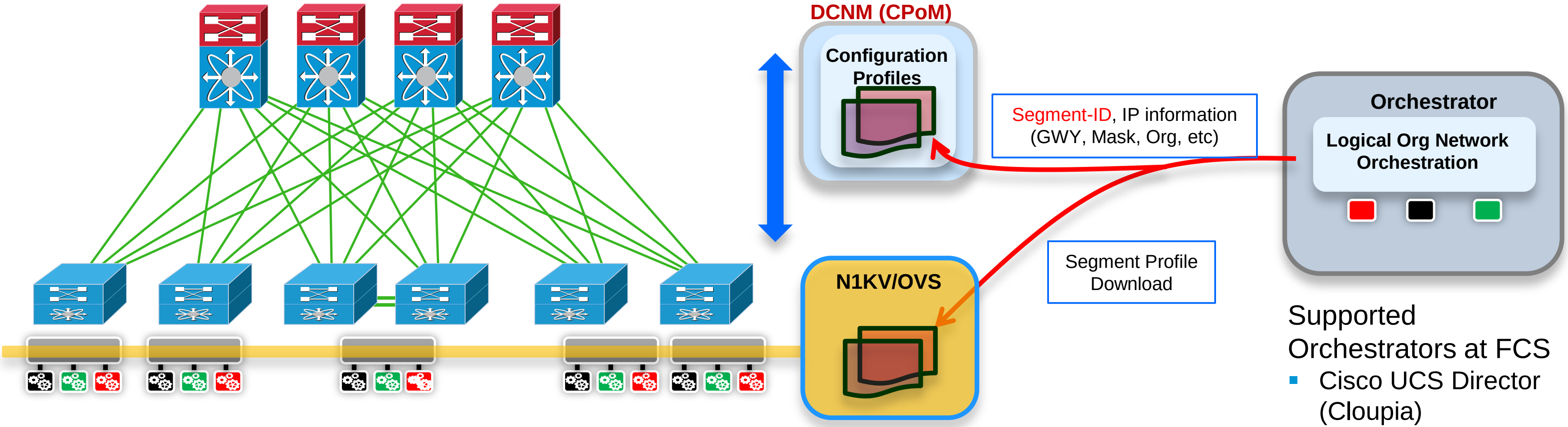


Workload Automation & Open Environment



*VDP (VSI Discovery and Configuration Protocol) is IEEE 802.1Qbg Clause 41

Workload Automation
Leveraging VDP for Leaf Auto-Configuration



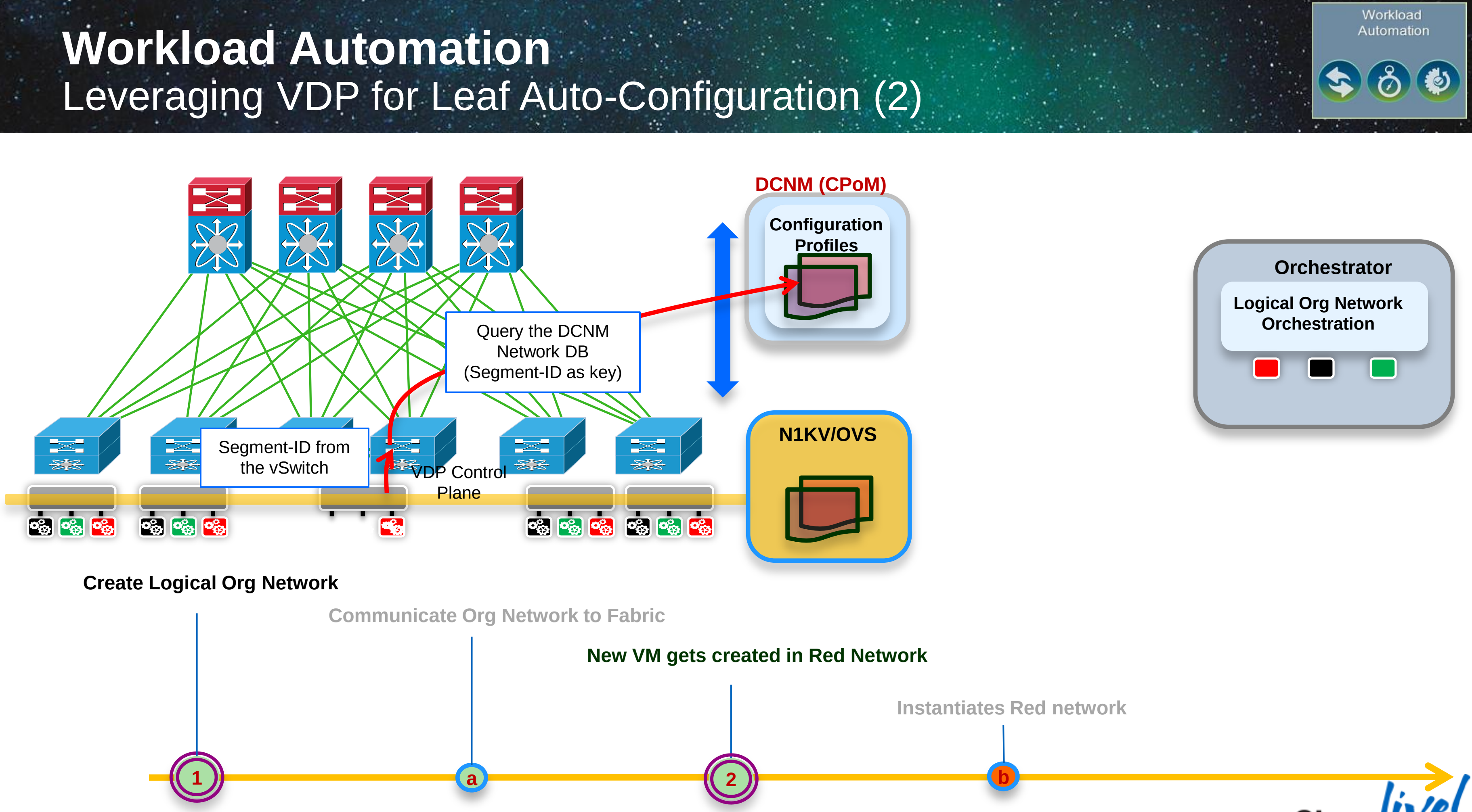
Supported Orchestrators at FCS
- Cisco UCS Director (Cloupia)
- OpenStack
- vCloud Director

Create Logical Org Network

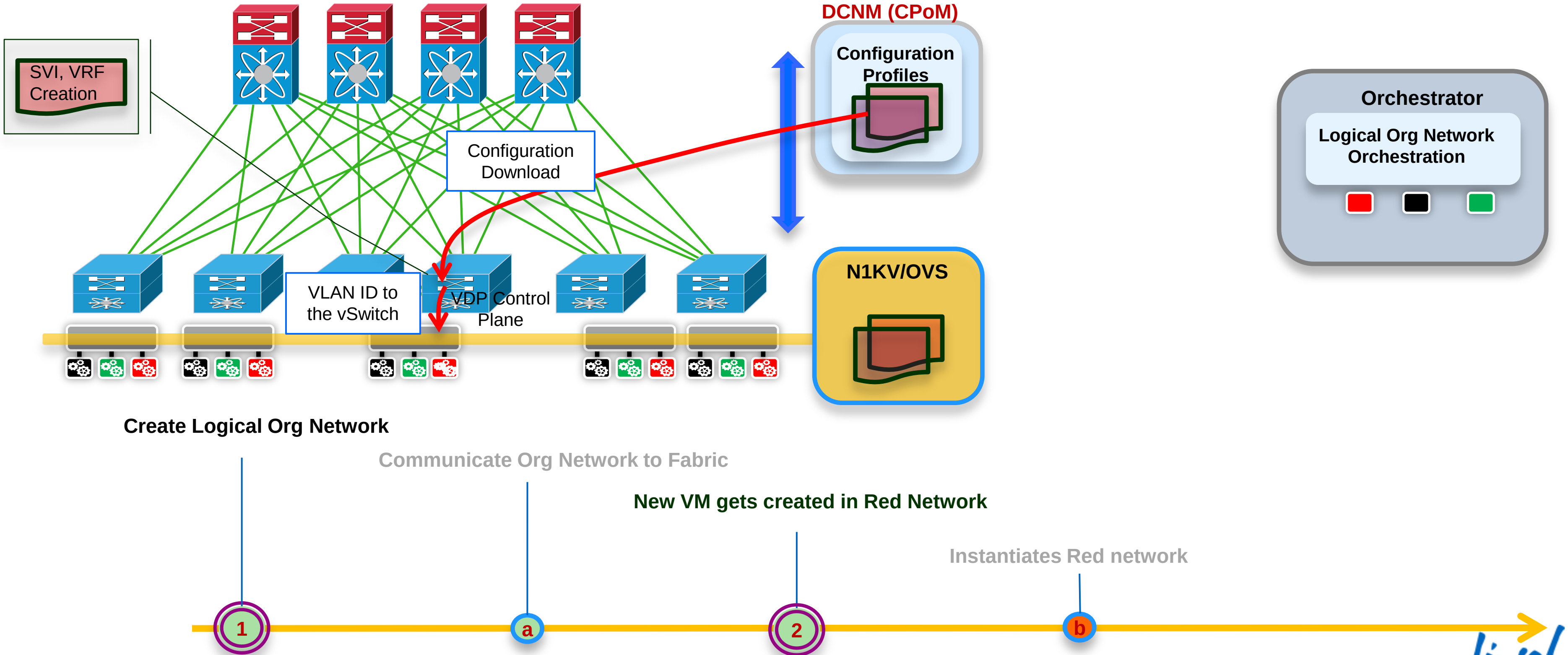
Communicate Org Network to Fabric

Workload Automation

Leveraging VDP for Leaf Auto-Configuration (2)



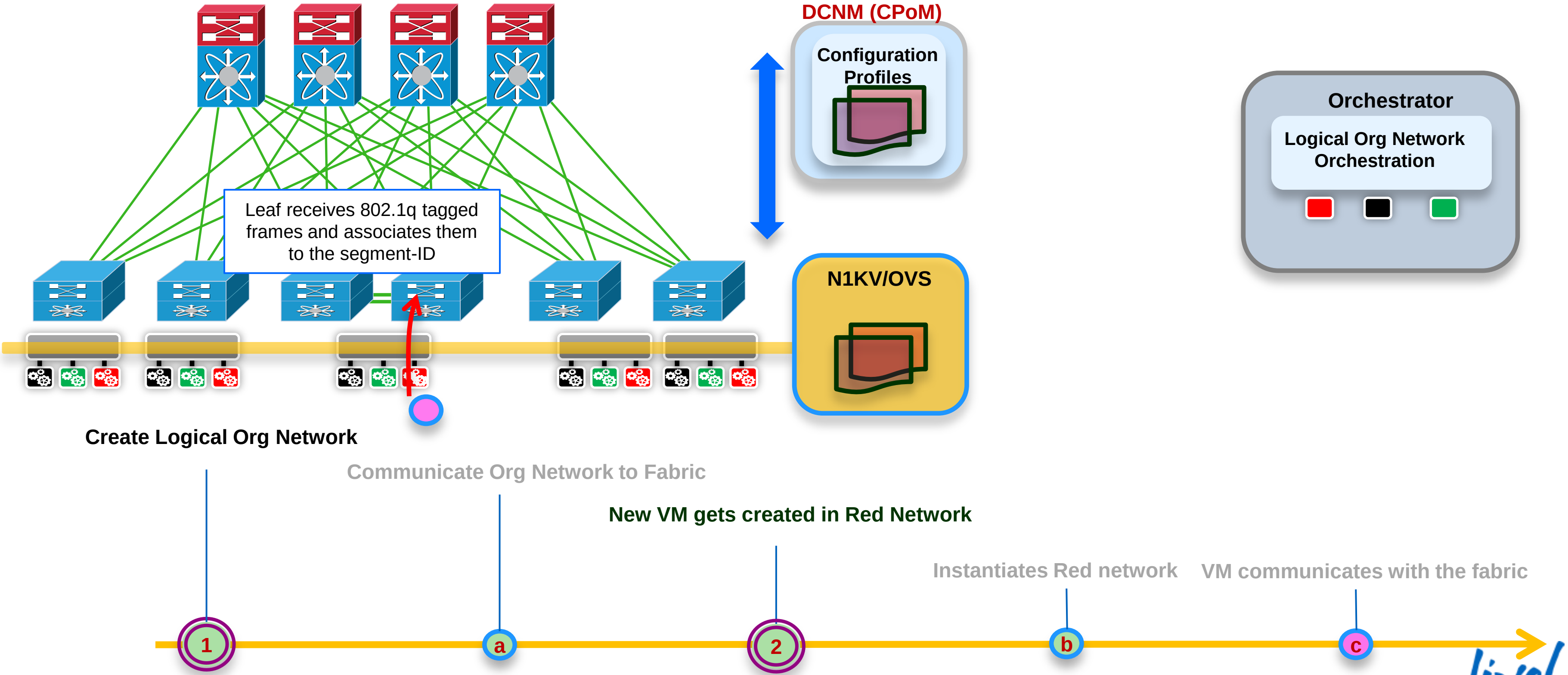
Workload Automation
Leveraging VDP for Leaf Auto-Configuration (3)



*VDP (VSI Discovery and Configuration Protocol is part of 802.1Qbg Draft

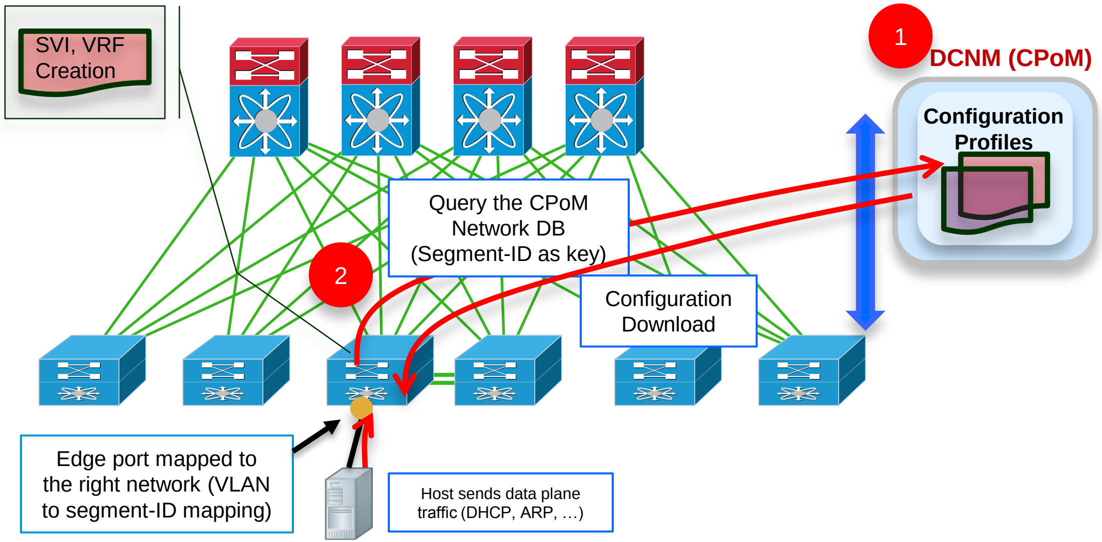
Workload Automation

Leveraging VDP for Leaf Auto-Configuration (4)



Workload Automation

What about Auto-Configuration for Physical Hosts?



Two steps required to provide connectivity into the fabric to a physical host

1. Adding configuration profile to the CPOM network database
2. Detecting when the host connects to query the database and instantiate the configuration on the leaf

Same model could apply to VMs deployed on vSwitches not supporting VDP

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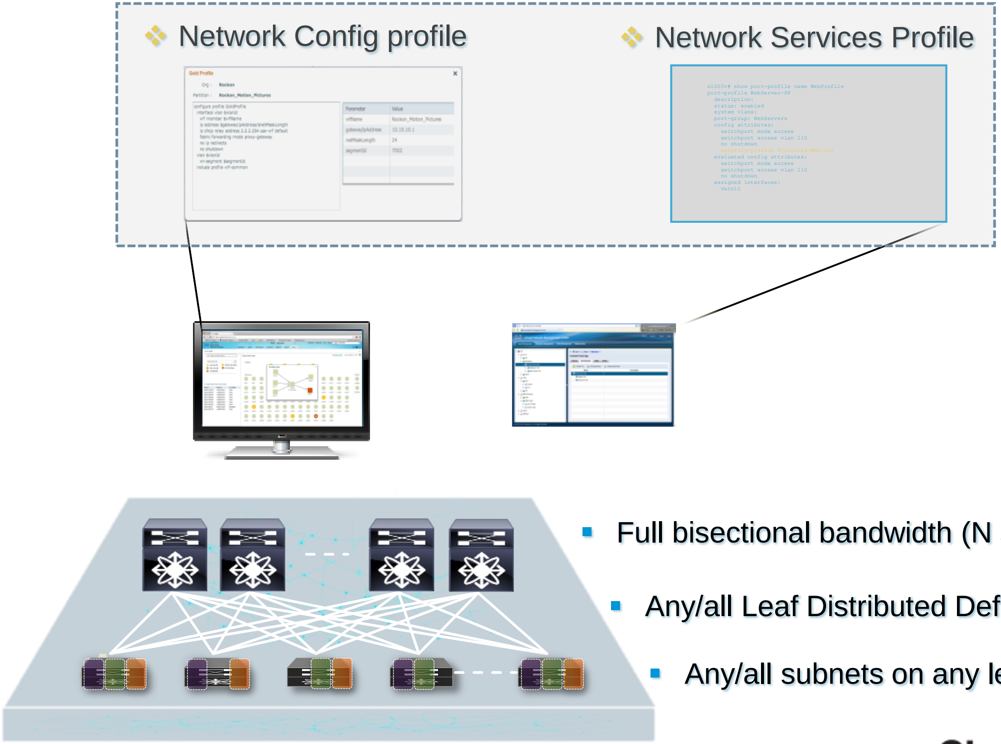
Cisco Dynamic Fabric Automation
Scale, Resiliency and Efficiency

Optimized Network

Advantages

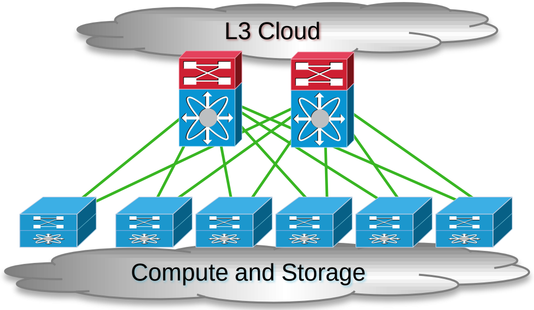
- Any subnet, anywhere, rapidly
- Reduced Failure Domains
- Extensible Scale & Resiliency
- Profile Controlled Configuration



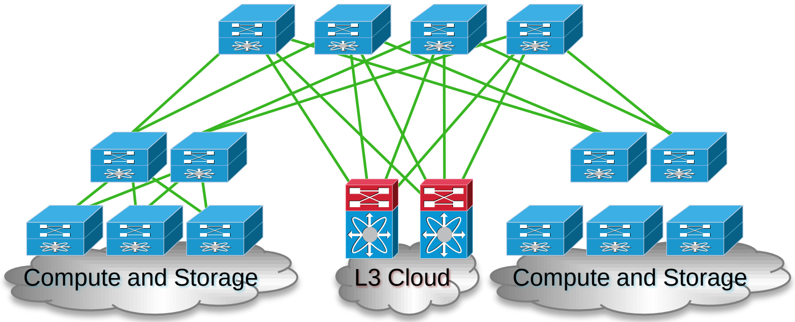
- Full bisectional bandwidth (N spines)
- Any/all Leaf Distributed Default Gateways
- Any/all subnets on any leaf

Cisco Dynamic Fabric Automation
Flexible Topologies Support

Traditional Access/Aggregation



Three Tiers (Fat Tree)



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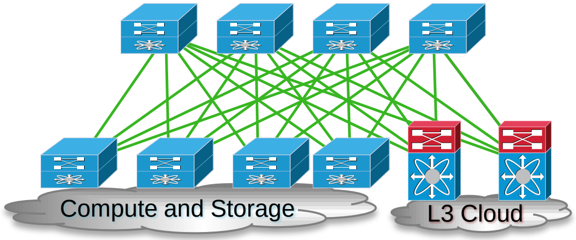
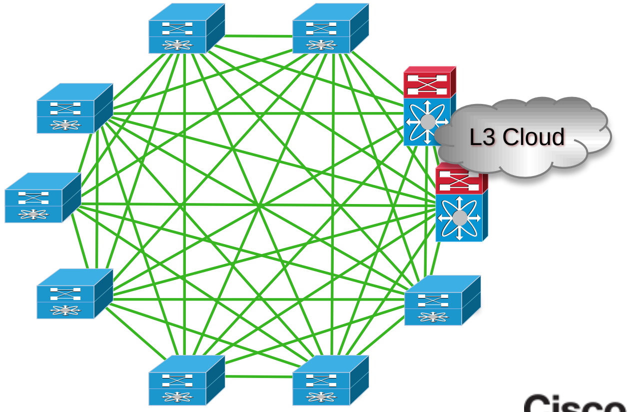


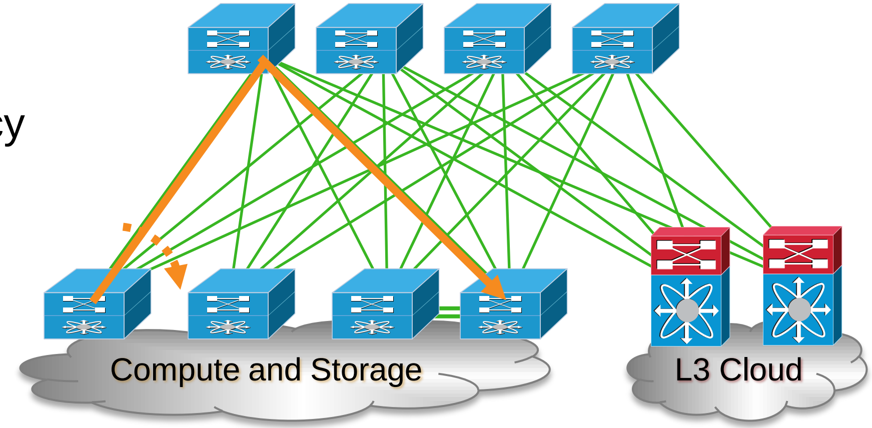
Full Mesh



Cisco Dynamic Fabric Automation
Fabric Properties

Optimized Network

- High Bisectional Bandwidth
- Wide ECMP: Unicast or Multicast
- Uniform Reachability, Deterministic Latency
- High Redundancy: Node/Link Failure
- Line rate, low latency, for all traffic

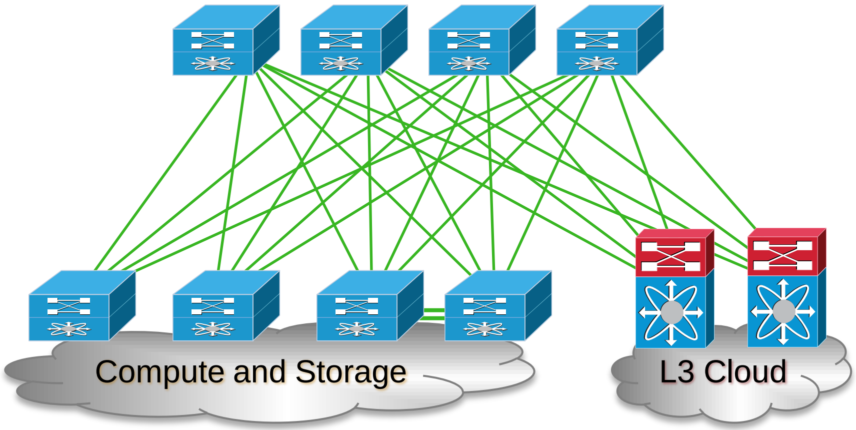


§§ Fabric properties applicable to all topologies §§

Cisco Dynamic Fabric Automation
Variety of Fabric Sizes

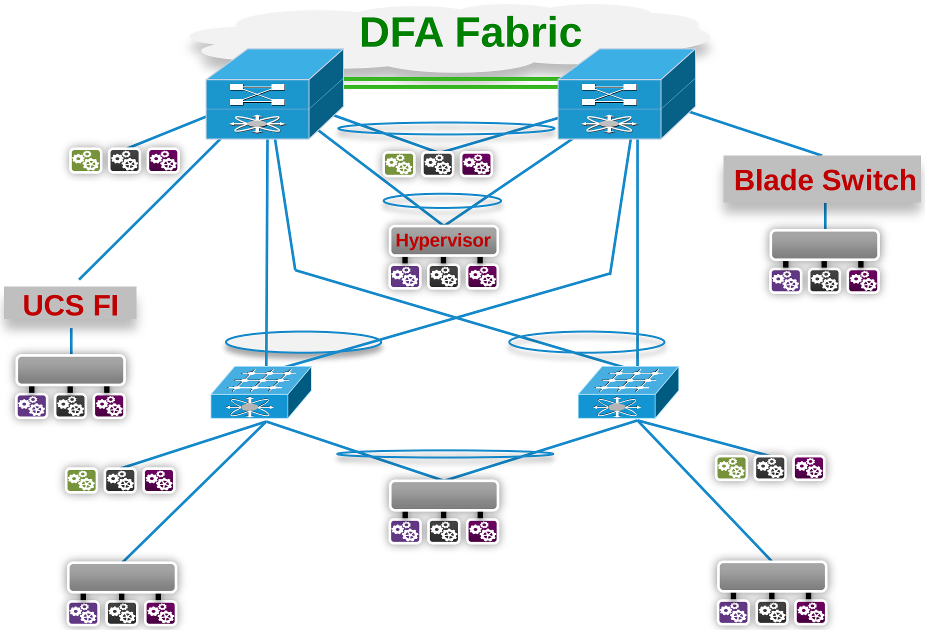
Optimized Network

- Fabric size: Hundreds to 10s of Thousands 10G ports
- Variety of Building Blocks:
  - ✓ Varying Size
  - ✓ Varying Capacity
  - ✓ Desired oversubscription
  - ✓ Modular and Fixed
- Scale Out Architecture
  - ✓ Add compute, service, external connectivity as the need grows



Cisco Dynamic Fabric Automation

Variety of South-bound Topological Connectivity



- Flexible connectivity options to the leaf nodes
  - ✓ FEX in straight-through or dual-active mode (eVPC)
  - ✓ UCS Fabric Interconnects
  - ✓ Hypervisors or bare-metal servers attached in vPC mode
- The FEX works as “remote linecards” and do not participate in DFA control plane and data plane encapsulation

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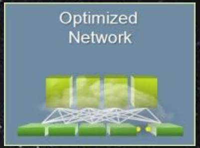
Control Plane
1 - IS-IS as Fabric Control Plane

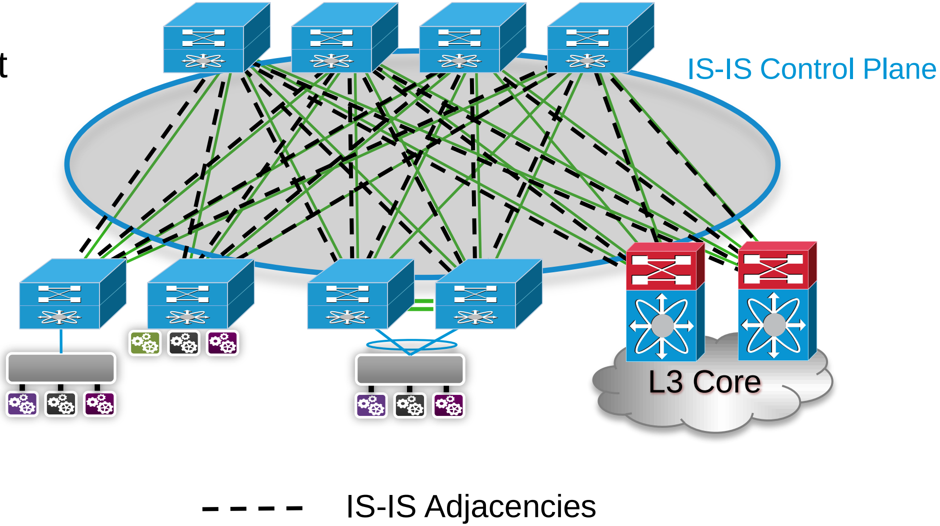


ISIS for fabric link state distribution

- Fabric node reachability for overlay encap
- Building multi-destination trees for multicast and broadcast traffic
- Quick reaction to fabric link/node failure
- Enhanced for mesh topologies

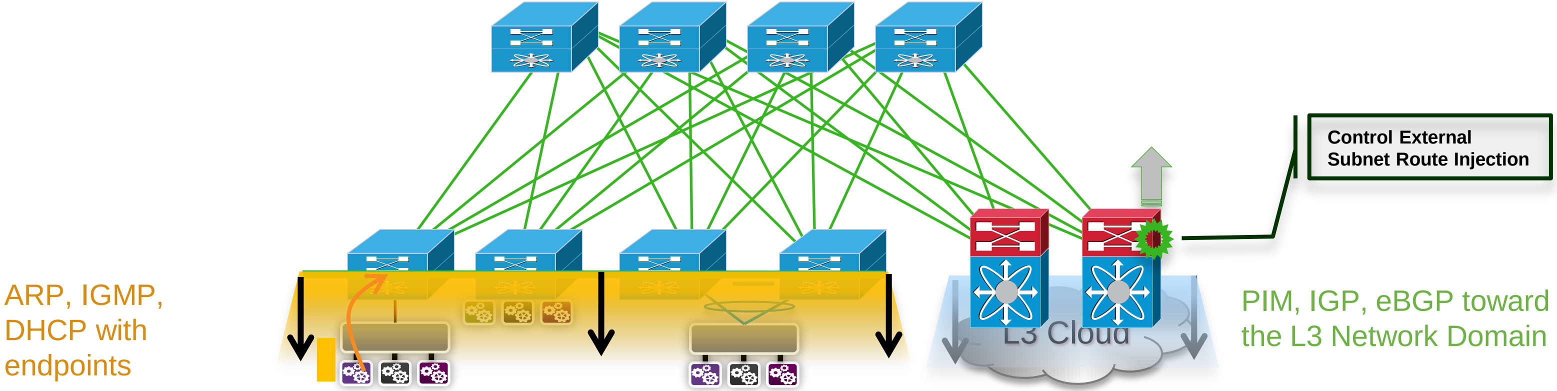
Fabric Control Protocol doesn't distribute

- Host Routes
- Host originated control traffic
- Server subnet information



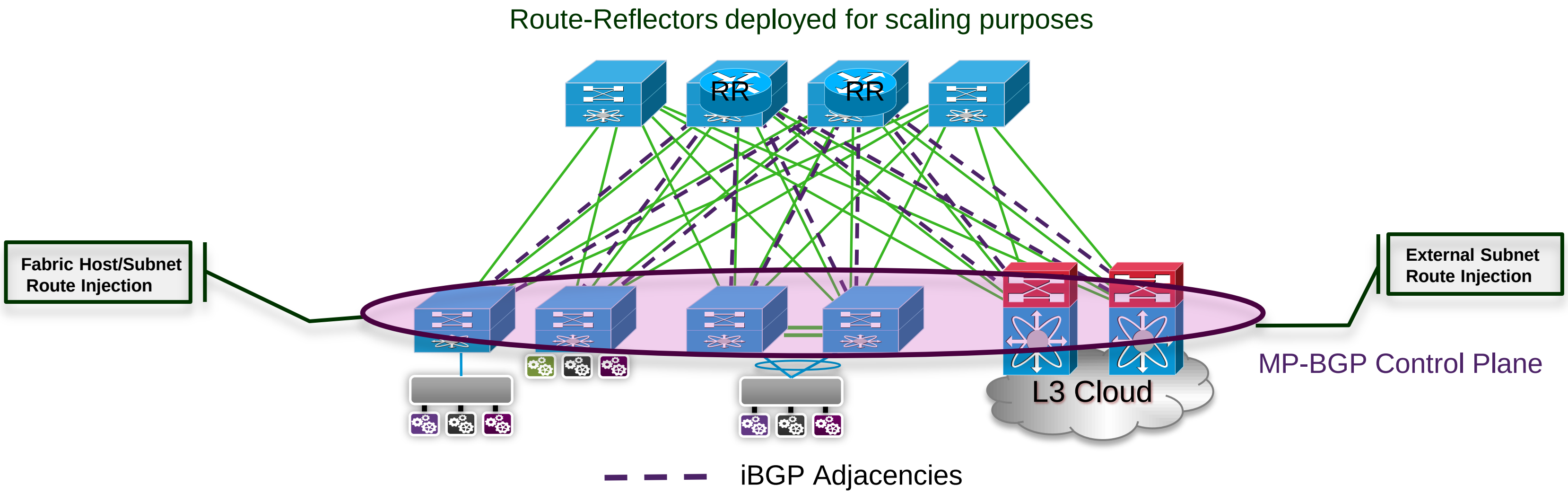
Control Plane

2 – Host Originated Protocols Containment



- ARP, IGMP, DHCP originated on servers are terminated on Leaf nodes
- Contain floods and failure domains, distribute control packet processing
- Terminate PIM, OSPF, eBGP from external networks on Border Leafs

Control Plane
3 – Host and Subnet Route Distribution



- Host Route Distribution decoupled from the Fabric link state protocol
- Use MP-BGP on the leaf nodes to distribute internal host/subnet routes and external reachability information
- MP-BGP enhancements to carry up to 100s of thousands of routes and reduce convergence time

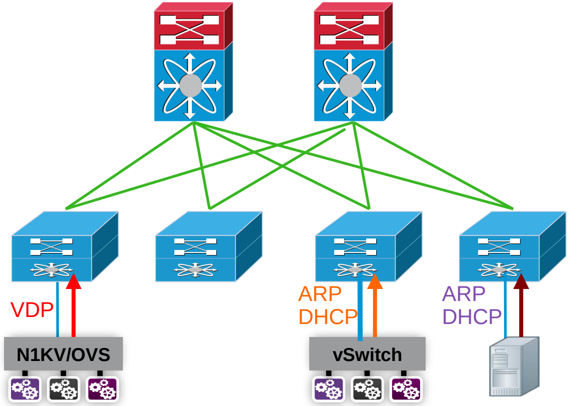
Control Plane
Hosts Detection and Deletion

Optimized Network

- In order to advertise host reachability information, a leaf must discover first locally connected devices
- Detection of local hosts

  Based on VDP or ARP/DHCP
- Detection of remote hosts

  Received MP-BGP notifications



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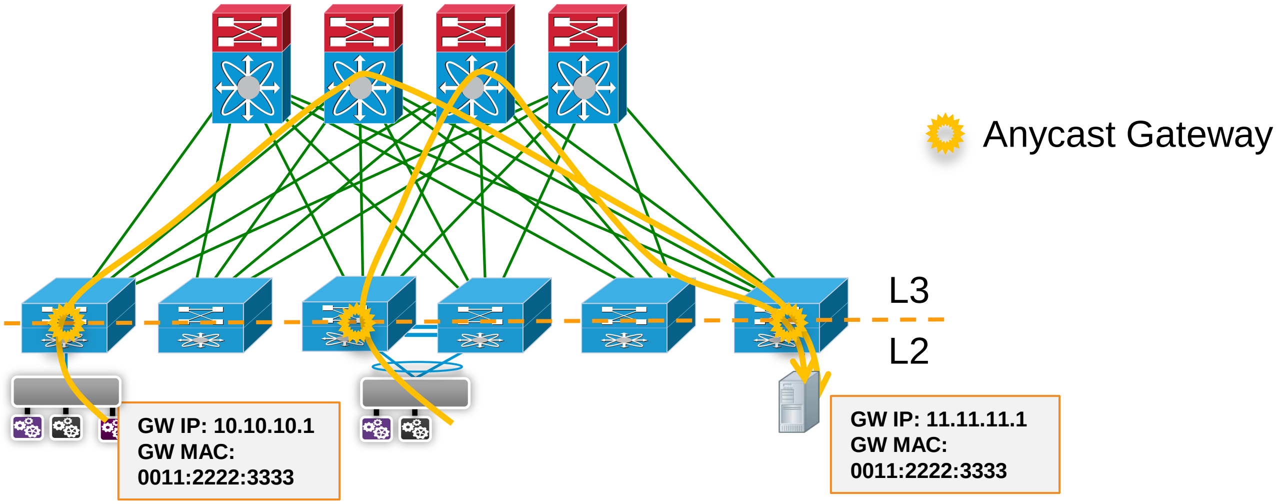
Optimized Network
Distributed Gateway at the Leaf



- Any subnet anywhere => Any leaf can instantiate any subnet

  All leafs share gateway IP and MAC for a subnet (No HSRP)

  ARPs are terminated on leafs, No Flooding beyond leaf

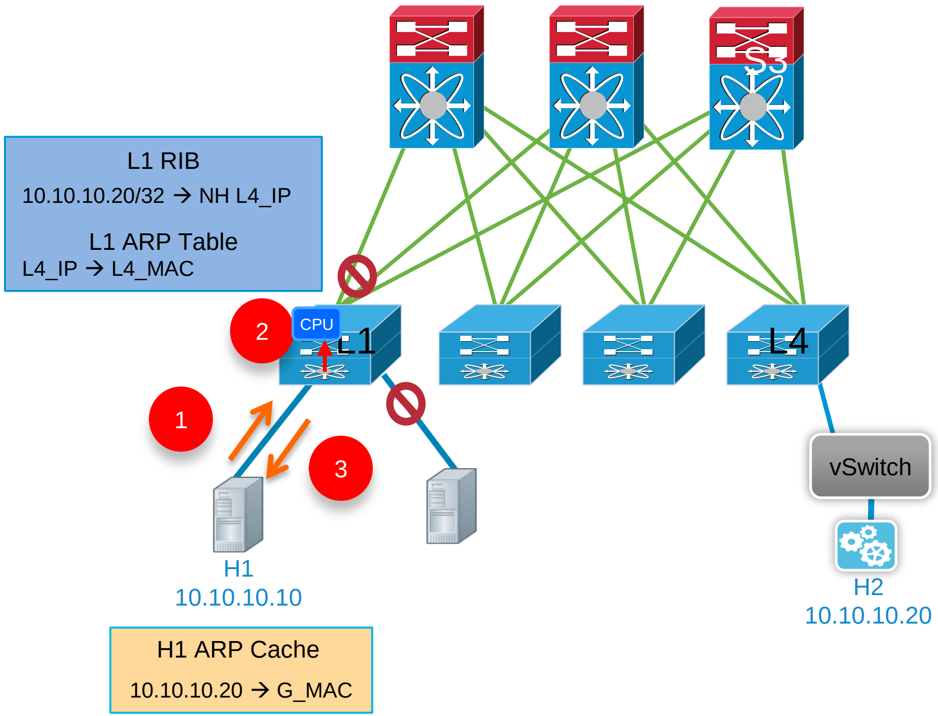
- Facilitates VM Mobility, workload distribution, arbitrary clustering
- Seamless L2 or L3 communication between physical hosts and virtual machines

Optimized Network

IP Forwarding within the Same Subnet

1. H1 sends an ARP request for H2 – 10.10.10.20
2. The ARP request is intercepted at the leaf L1 and punted to the Sup
3. Assuming a valid route to H2 does exist in the Unicast RIB, L1 sends the ARP reply with the G_MAC so that H1 can build its ARP cache

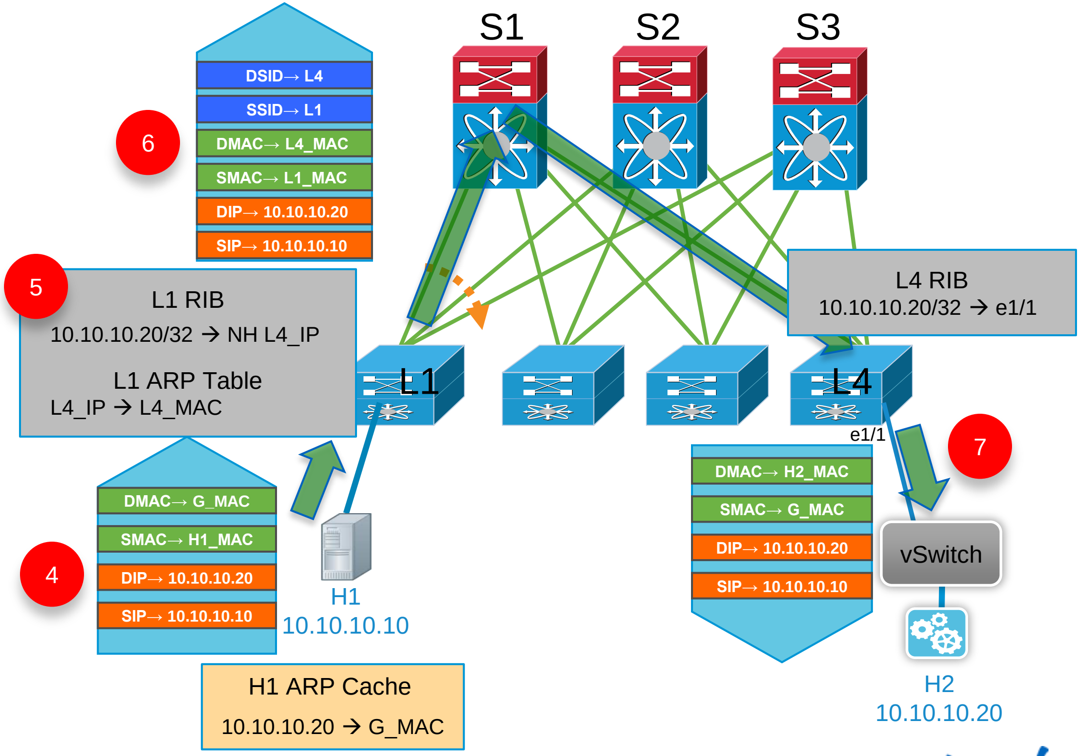
Note: the ARP request is NOT flooded across the Fabric, nor out of other local interfaces belonging to the same L2 domain



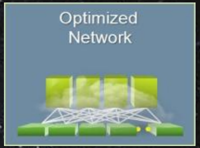
Optimized Network

IP Forwarding within the Same Subnet (2)

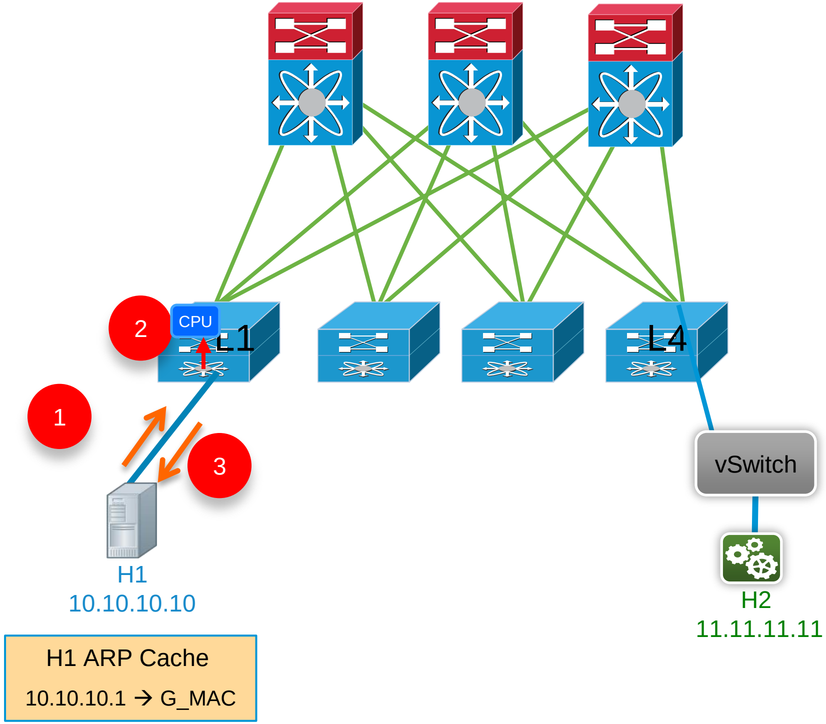
4. H1 generates a data packet with G_MAC as destination MAC
5. L1 receives the packet, remove the L2 header and performs Layer 3 lookup for the destination
6. L1 adds the Layer 2 and the FP headers and forwards the FP frame across the Fabric, picking one of the 3 equal cost paths available via S1, S2 and S3
7. L4 receives the packet, strips off the FP and L2 headers and performs L3 lookup and forwarding toward H2



Optimized Network
IP Forwarding Across Different Subnets



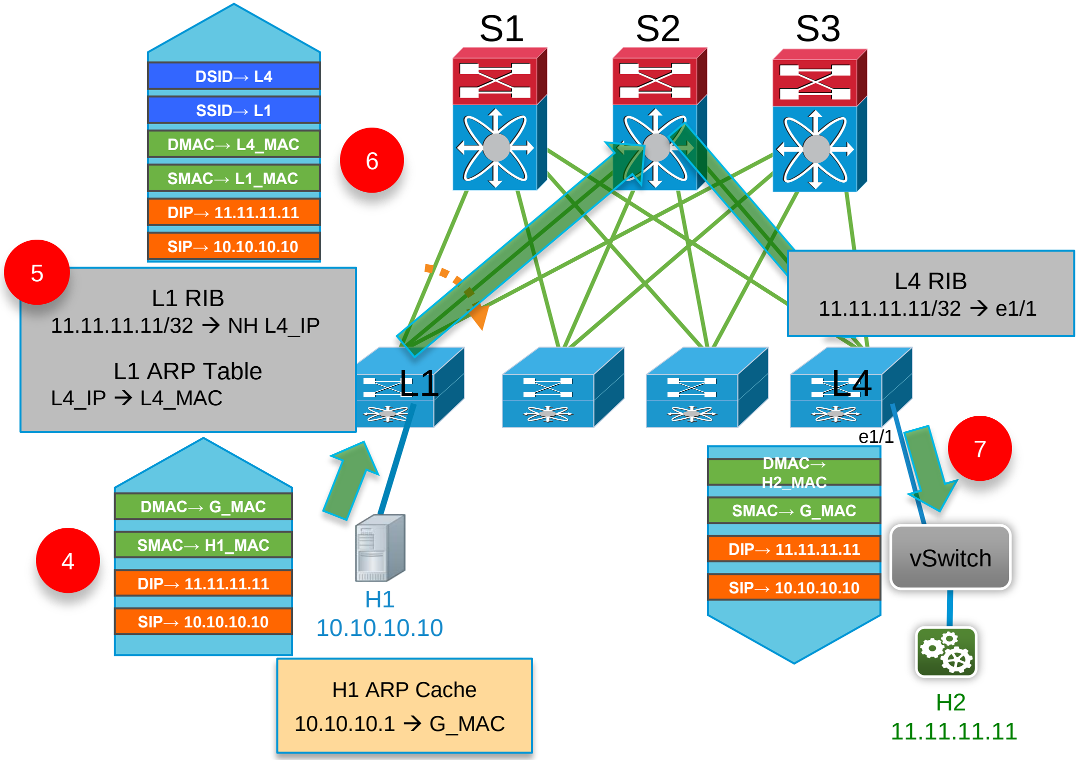
1. H1 sends ARP request for default gateway – 10.10.10.1
2. The ARP request is intercepted at the leaf and punted to the Sup
3. L1 acts as regular default gateway and sends ARP reply with G_MAC



Optimized Network

IP Forwarding Across Different Subnets (2)

4. H1 generates a data packet destined to H2 IP with G_MAC as destination MAC
5. L1 receives the packet, remove the L2 header and performs Layer 3 lookup for the destination
6. If valid routing information for H2 are available in the unicast routing table, L1 adds the Layer 2 and the FP headers and forwards the FP frame across the Fabric, picking one of the 3 equal cost paths available via S1, S2 and S3
7. L4 receives the packet, strips off the FP and L2 headers and performs L3 lookup and forwarding toward H2



Optimized Network
Introducing L3 Conversational Learning

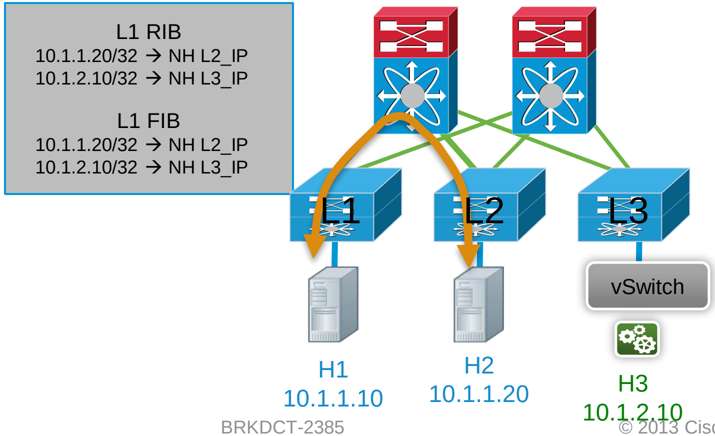
- Use of /32 host routes may lead to scaling issues if all the routes are installed in the HW tables of all leaf nodes

    L3 conversational learning is introduced to alleviate this concern

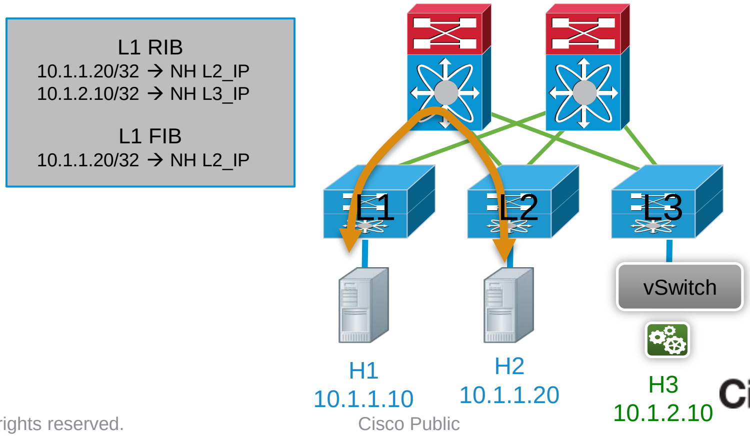
    Disabled by default → all host routes are programmed in the HW

- With L3 conversational learning, host routes for remote endpoints will be programmed into the HW FIB (from the SW RIB) upon detection of an active conversation with a local endpoint

Default Behavior (No L3 Conversational Learning)



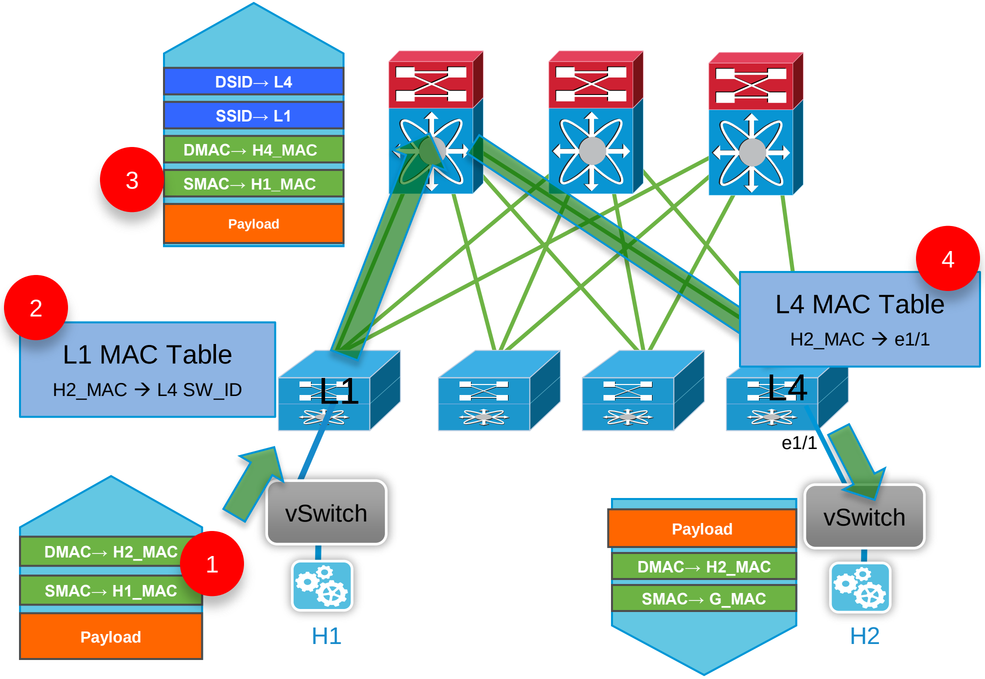
After Enabling L3 Conversational Learning



Forwarding
L2 non IP Flows

Optimized Network

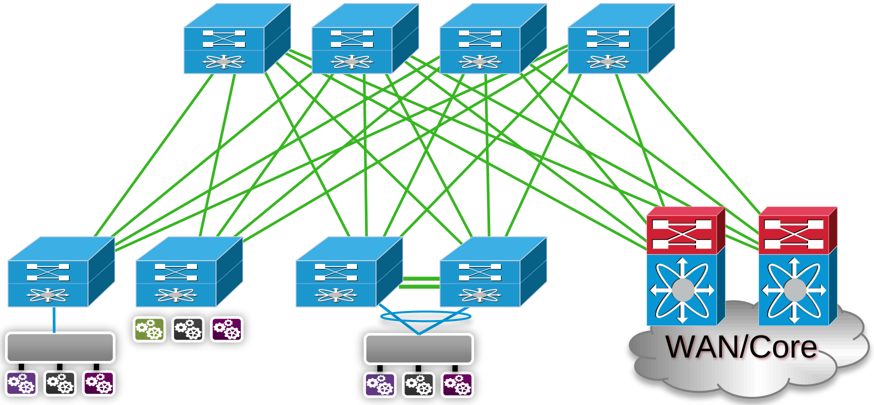
1. H1 originates a packet destined to H2 MAC address
2. L2 lookup is performed by L1 in the MAC Table for the VLAN the frame belongs to
3. L1 adds the Layer 2 and FP headers before sending the packet into the fabric
4. L4 receives the frame, decapsulates the FP header, performs the L2 lookup and then sends it to H2



Optimized Network
Multicast Forwarding

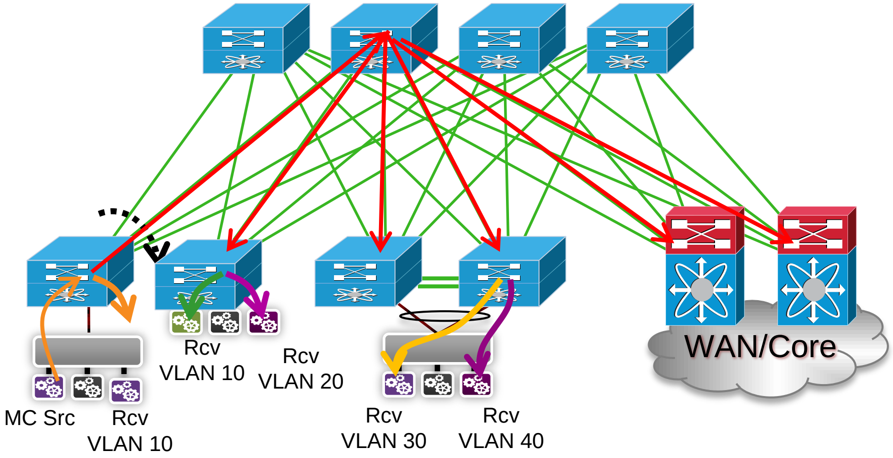
Optimized Network

- Fabric supports computation of multiple distribution trees leveraging IS-IS
    - Used for multicast and broadcast traffic
    - No need for other multicast protocols (PIM, etc.) inside the fabric
- Multi Destination Trees (MDTs) Rooted on Spines
- Ingress Leaf load balances traffic across multiple paths
    - Efficient use of fabric links



Optimized Network
Multicast Forwarding

- Two tiers multicast replication across the fabric Ingress
  - Ingress Leaf always performs multicast routing functions and sends a single copy onto the fabric
  - Spine node replicates to the leaf nodes
  - Destination Leaf nodes locally replicate to server ports across subnets
- Optimization possible to allow pruning on the spine (per tenant/VRF or per group)



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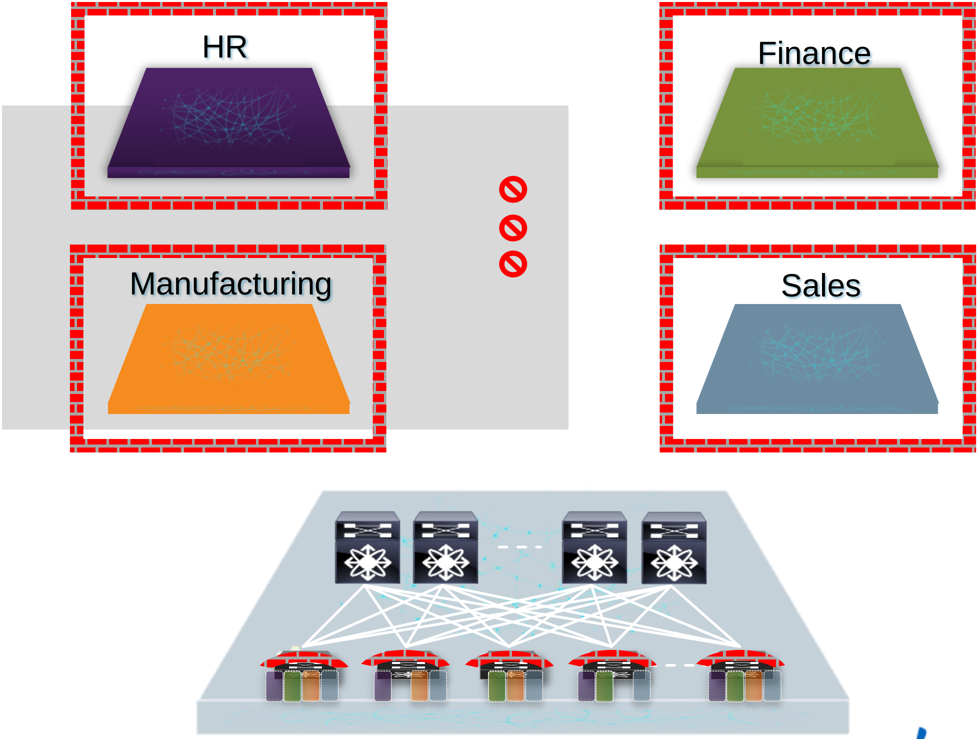
Virtual Fabrics for Public or Private Cloud Environments

Advantages

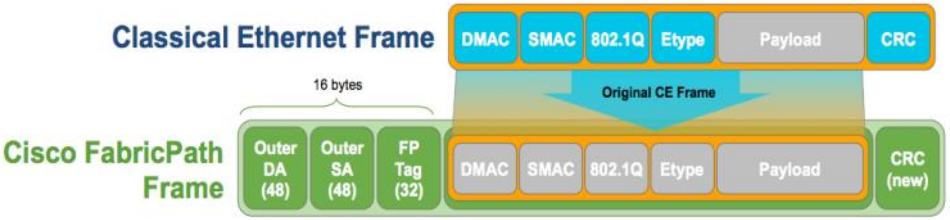
- Any workload, any vFabric, rapidly
- Scalable Secure vFabrics
- vFabric Tenant Visibility
- Routing/Switching Segmentation



Virtual Fabrics
Introducing Segment-ID Support

- Traditionally VLAN space is expressed over 12 bits (802.1Q tag)
  - Limits the maximum number of segments in a datacenter to 4096 VLANs

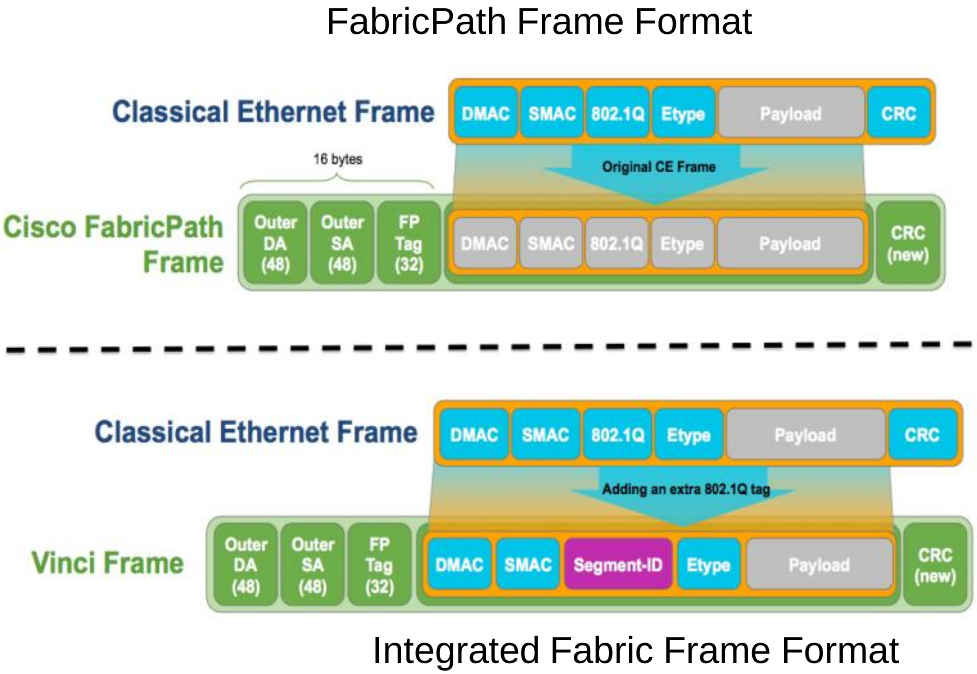
FabricPath Frame Format



Virtual Fabrics

Introducing Segment-ID Support

- Traditionally VLAN space is expressed over 12 bits (802.1Q tag)
    - Limits the maximum number of segments in a datacenter to 4096 VLANs
- DFA leverages a double 802.1Q tag for a total address space of 24 bits
    - Support of ~16M L2 segment (10K targeted at FCS)
- Segment-ID is hardware-based innovation offered by leaf and spine nodes part of the Integrated Fabric

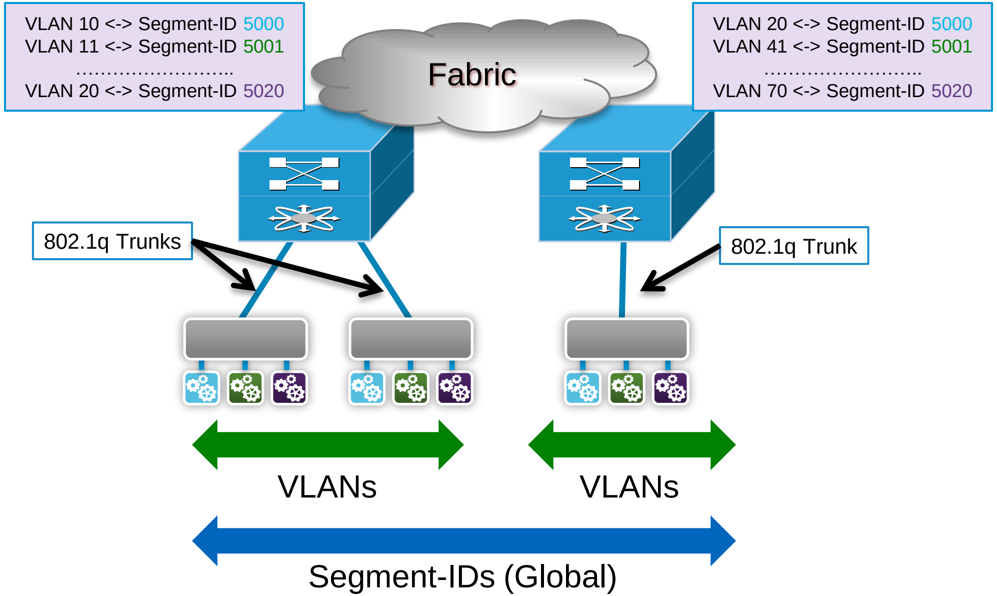


Virtual Fabrics

802.1Q Tagged Traffic to Segment-ID Mapping

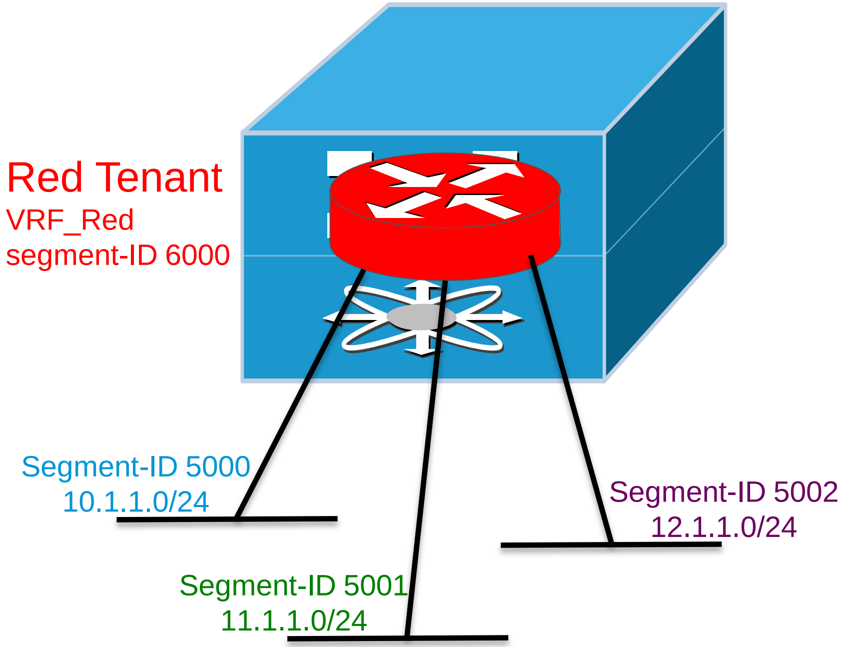
- Segment-IDs are utilized for providing isolation at L2 and L3 across the Integrated Fabric
- 802.1Q tagged frames received at the leaf nodes from edge devices must be mapped to specific Segments
- The VLAN-Segment mapping can be performed on a leaf device level

  VLANs become locally significant on the leaf node and 1:1 mapped to a Segment-ID
- Segment-IDs are globally significant, VLAN IDs are locally significant



Virtual Fabrics
How are Segment-IDs Utilized?

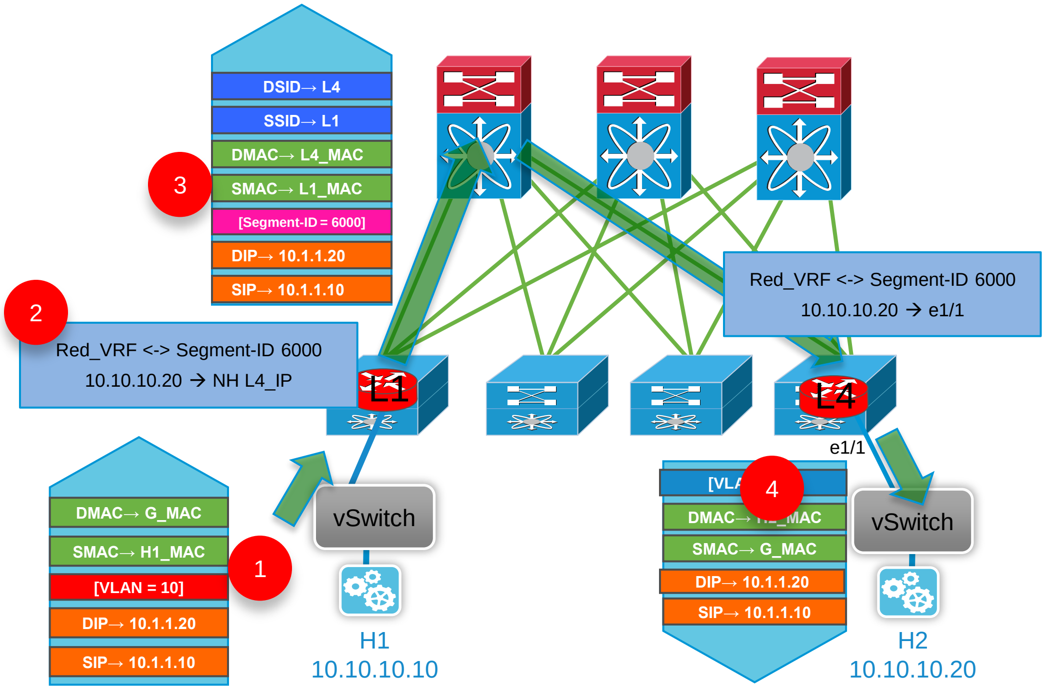
- Each IP subnets defined at the edge of the Integrated Fabric is associated to a Layer 2 domain, which is represented by a Segment-ID
- Multiple Segments can be defined for a given tenant and are usually mapped to a L3 VRF uniquely identifying that tenant
- A dedicated Segment-ID value uniquely identifies each VRF defined in the Integrated Fabric



Virtual Fabrics
Fabric Routed Flows

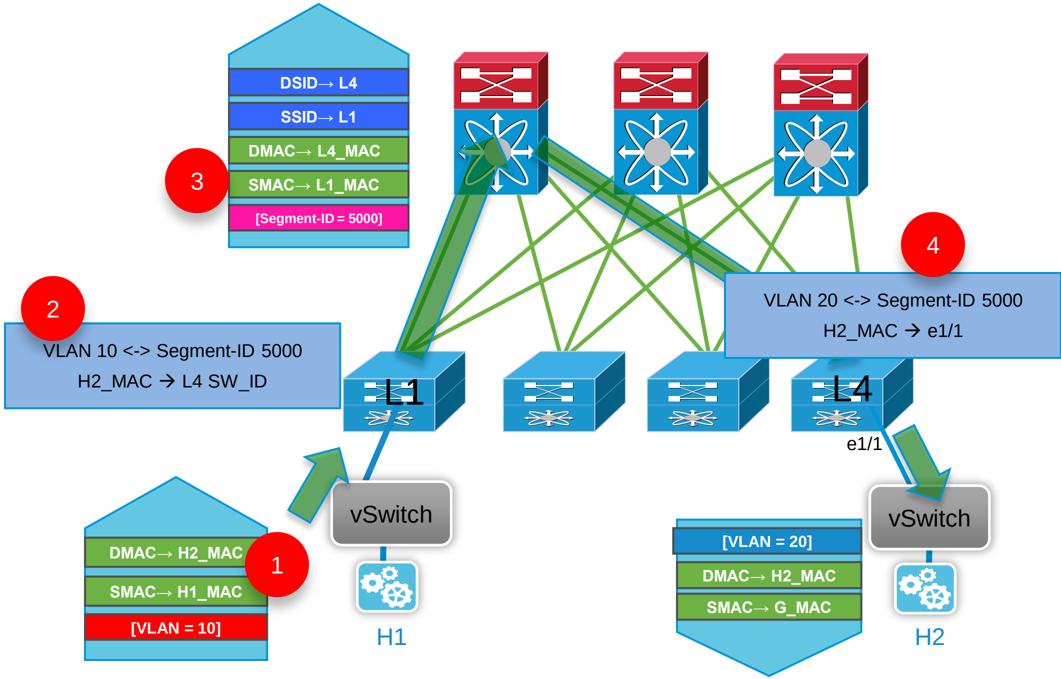
1. H1 sends a packet to H2 → traffic between the vSwitch and the Leaf is tagged with a local VLAN-ID 10
2. L3 lookup is performed by L1 in the context of the Red VRF
3. L1 adds the L2 and FP headers before sending the packet into the fabric. The Segment-ID identifying the Red VRF is added inside the L2 header
4. L4 receives the frame and associates it to the Red VRF by looking at the Segment-ID value. It then sends it to H2 using a local VLAN-ID 20

Note: this behavior applies to all fabric routed flows (intra-subnet or inter-subnet)



Virtual Fabrics
L2 non IP Flows

1. H1 sends a packet to H2 → traffic between the vSwitch and the Leaf is tagged with a local VLAN-ID 10
2. L2 lookup is performed by L1 in the MAC Table for the Segment-ID associated to VLAN 10 (5000)
3. L1 adds the L2 and FP headers before sending the packet into the fabric. The Segment-ID associated to VLAN 10 is added inside the L2 header
4. L4 receives the frame and performs the L2 lookup by looking at the Segment-ID value. It then sends it to H2 using a local VLAN-ID 20



Agenda

- DFA Requirements and Functions
- Fabric Management
- Workload Automation
- Optimized Network
  - Fabric Properties
  - Control Plane
  - Forwarding Plane
- Virtual Fabrics
- Hardware Support

Cisco Dynamic Fabric Automation

Platform Support at FCS

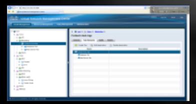
Cloud Stacks & Orchestration Tools

Compute & Storage

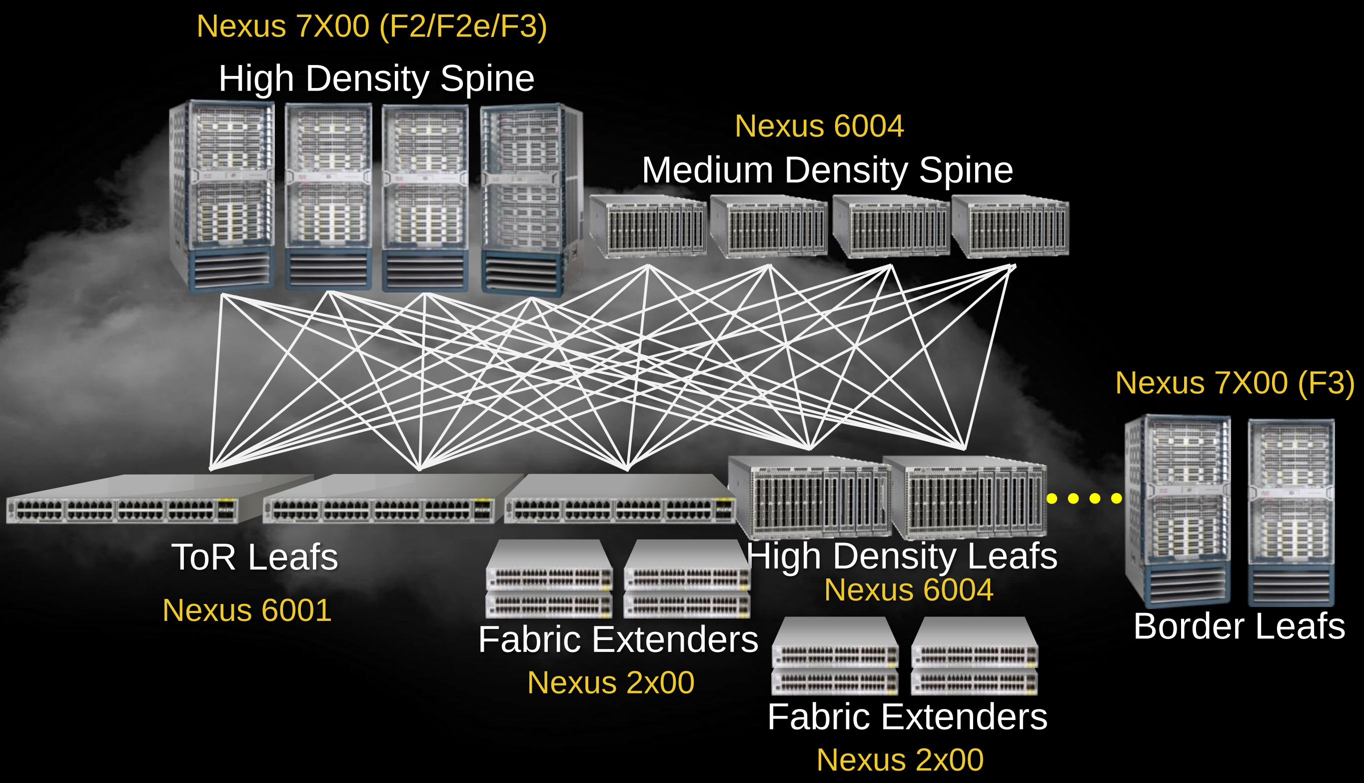
Network

DCNM/CPoM

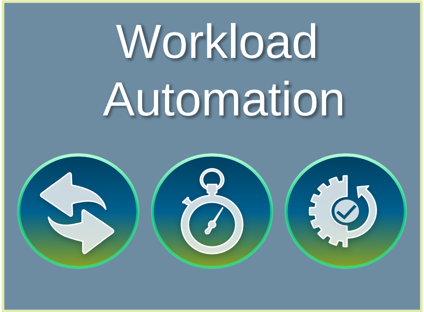
Network Services

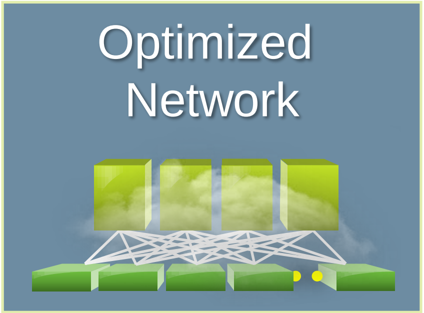


Services Controller



Cisco Dynamic Fabric Automation Architecture
Where to Get More Information





TECDCT-2306 → Thu June 27th 8 am-12 pm

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